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Functoriality of inference on diagrams in the category of Markov kernels.

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Inference on diagrams

Introduction to geometric deep learning [BBCV21]:

- Deep learning ← curse of dimensionality
- Accounting for symmetry
 - \rightarrow Translation \rightsquigarrow CNN
 - \rightarrow Other groups [SPMBO22]
- Geometry ~> discretize
 - → Graph NN [BBCV21]
 - ightarrow Nodes share same features
 - \rightarrow Limitations: heterogeneous data
- Heterogeneity
 - \rightarrow Cellular sheaves [Cur13]
 - $\rightarrow~$ Cell complex, faces \rightsquigarrow feature space, inclusions \rightsquigarrow linear maps
 - \rightarrow Functor from a poset to Vect [SP22, SP24a, SPR24, SP24b]
 - → Sheaf Neural Networks [BGC⁺22]

- Our focus:
 - $\rightarrow\,$ Bayesian inference, graphical models, Markov random fields, factor graphs

Context:

Let *X* be a random variable taking values in a finite measurable space E_X , and $\theta \in \Theta \subseteq \mathbb{R}^d$ that parametrizes a collection of probability measures $P(x|\theta)$, where $x \in E_X$.

Assume that one is given a prior $Q \in \mathbb{P}(\Theta)$, where $\mathbb{P}(\Theta)$ denotes the space of probability measures on Θ . For any observation $x_0 \in E_X$, one computes the posterior using Bayes' rule:

$$P(\theta|x) = \frac{P(x|\theta)Q(\theta)}{\int P(x|\theta)Q(\theta) \, d\theta}$$

Exponential Computational Cost

- <u>Problem</u>: When $\theta = (\theta_i, i \in [0, N])$ is a collection of variables, where each $\theta_i \in E$.
- $\sum_{\theta} P(x|\theta)Q(\theta) = \sum_{\theta_0} \cdots \sum_{\theta_N} P(x|\theta_0, \dots, \theta_N)Q(\theta_0, \dots, \theta_N)$

 \rightarrow Number of operations: $\mathscr{O}(|E|^N)$

- <u>Notation</u>: For a set of indices *I* and a subset $a \subseteq I$, $\theta_a := (\theta_i \in E_i, i \in a)$.
- In what follows, all the sets in which variables take values will be finite: *E_i* are finite sets.

• One relation to statistical mechanics:

Let $\theta = (Y_i \in E_i, i \in I)$ be the unobserved variables, and $X = (X_j \in F_j, j \in J)$ the observed variables. Both *I* and *J* are finite sets, and E_i for $j \in J$ and F_i for $i \in I$ are finite sets.

$$\ln P(\theta, X) = -\beta \sum_{a \subseteq I \sqcup J} H_a(X_{a \cap J}, Y_{a \cap I})$$

Given an observation $x = (x_i, i \in I)$, computing $\ln P(\theta|x)$ is equivalent to computing:

$$\ln \sum_{(y_i, i \in I)} e^{-\beta \sum_{a \subseteq I \sqcup J} H_a(x_{a \cap J}, y_{a \cap I})}$$

This is the same as:

$$\ln Z(x) := \ln \sum_{\theta \in \prod_i E_i} e^{-eta \widetilde{H}_x(heta)}$$

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• Similar frameworks but different names: Bayesian networks, graphical models, factor graphs, Markov random fields.

Definition (Factorisation Space)

Let *I* be a finite set, and let $\mathscr{A} \subseteq \mathscr{P}(I)$, where $\mathscr{P}(I)$ is the power set of *I*. Let $(E_i, i \in I)$ be a collection of sets, and let $E_a = \prod_{i \in a} E_i$ for any $a \in \mathscr{P}(I)$. For $x \in \Omega$, we denote by x_a its projection onto E_a . The factorisation space over \mathscr{A} is defined as follows:

$$\mathsf{Fac}_{\mathscr{A}} = \{ P \in \mathbb{P}(\Omega) : \exists (f_a \in \mathbb{R}_{>0}^{E_a}, a \in \mathscr{A}) \text{ s.t. } \forall x \in \Omega, \ P = \prod_{a \in \mathscr{A}} f_a(x_a) \}$$

Consider an undirected graph G = (I, A) that is **acyclic**. Denote $\mathscr{A}(G)$ as the *partially ordered set* (poset) with elements $V = I \sqcup A$ and the following relations:

- $\forall i \in I, i \leq i$, and $\forall e \in A, e \leq e$
- $\forall i \in I, \ \forall e \in A, \ i \leq e \iff i \in e$

Proposition (Factorization on Acyclic Graphs)

Let I be a finite set, and let $\Omega = \prod_{i \in I} E_i$ be a product of finite sets, and $X_i, i \in I$, a collection of random variables taking values respectively in E_i . Let G = (I, A) be a finite acyclic graph. $P_X \in \mathbb{P}_{>0}(E)$ factors according to $\mathscr{A}(G)$, i.e., $P_X \in \operatorname{Fac}_{\mathscr{A}(G)}$, if and only if for any $\omega \in \Omega$,

$$P_X(\omega) = \frac{\prod_{e \in A} P_{X_e}(\omega_e)}{\prod_{i \in I} P_{X_i}^{d(i)-1}(\omega_i)},$$

where d(i) is the degree of node $i \in I$.

- Bayesian inference is maximizing (relative) entropy.
- Entropy:

$$S(Q) = -\sum_{\omega \in E} Q(\omega) \ln Q(\omega)$$
 (0.1)

 Recall that minimizing Gibbs free energy gives Helmholtz free energy:

$$\beta \frac{-\ln Z}{\beta} = \inf_{Q \in \mathbb{P}(E)} (\mathbb{E}_Q[\beta H] - S(Q))$$

• Set $\beta = 1$.

• But entropy:

$$\mathcal{S}(\mathcal{P}_X) = \sum_{e \in \mathcal{A}} \mathcal{S}(\mathcal{P}_{X_e}) - \sum_{i \in I} (d(i) - 1) \mathcal{S}(\mathcal{P}_{X_i})$$

- Inclusion-exclusion formula: c(e) = 1, c(i) = -(d(i) 1)
- Remarkably, Bayesian inference is the same as minimizing [YFW05, YFW03]:

$$\mathcal{F}_{\mathsf{Bethe}}(\mathcal{Q}) = \sum_{a \in V} c(a) \left(\mathbb{E}_{Q_a}[\mathcal{H}_a] - S(Q_a)
ight)$$

where $Q := (Q_a \in \mathbb{P}(X_a), a \in V)$ with compatibility by marginalization:

 \rightarrow If *a* is an edge and *i* a vertex in *a*

$$\rightarrow \pi_i^e : E_e \rightarrow E_i$$

 \rightarrow We ask that $\pi_{i}^{e}(Q_{e}) = Q_{i}$

- Bayesian inference corresponds to computing $\ln Z$ for a Hamiltonian $H : \prod_{i \in I} E_i \to \mathbb{R}$, with $Z = \sum_x e^{-\beta H(x)}$.
- From now on, set $\beta = 1$; notation $E = \prod_{i \in I} E_i$.
- It is computationally costly to compute directly, but note that

$$-\ln Z = \inf_{Q\in\mathbb{P}(E)} (\mathbb{E}_Q[H] - S(Q))$$

• The previous problem can be reformulated as minimizing:

$$F_{ ext{Bethe}}(Q_a, a \in \mathscr{A}(G)) = \sum_{a \in \mathscr{A}(G)} c(a) \left(\mathbb{E}_{Q_a}[H_a] - S(Q_a)
ight)$$

with $Q_i(x_i) = \sum_{y \in X_{i'}} Q_e(x_i, y)$ when $e = \{i, i'\}$.

Belief propagation is an algorithm of complexity *O*(|*A*||*E_i*|²) to solve this optimization problem, when *E_i* = *E_j* for all *i*, *j* ∈ *I*.

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- \rightarrow Extension to higher-order interactions: not just graphs.
- \rightarrow I did not invent it [Pel20, YFW05]... but no name?

Definition (Graphical Presheaves)

Let *I* be a finite set and $\mathscr{A} \subseteq \mathscr{P}(I)$ be a sub-poset of the powerset of *I*. Let $E_i, i \in I$ be finite sets. For $a \in \mathscr{A}$, define $E_a := \prod_{i \in a} E_i$. Let $F(a) := E_a$, and for $b \subseteq a$, let $F_b^a : E_a \to E_b$ be the projection map from $\prod_{i \in a} E_i$ to $\prod_{i \in b} E_i$. We call *F* a graphical presheaf from \mathscr{A} to **Mes**^{*f*}.

- Only projections.
- Only products of variables, and subcollections of variables.

- Consider any map, not just projections:
 - \rightarrow Any measurable maps for $b \rightarrow a$ and even Markov kernels, i.e., stochastic matrices when the source and target are finite sets.
- Account for possible heterogeneity, incompleteness, and incompatibility in the description of variables:
 - \rightarrow Agents with different world models that communicate their beliefs.
 - ightarrow Broader class of effective potentials in computational chemistry.

Extension done in previous work [SP22, SPR24, SP24a]

- **Kern**^{*f*}: objects are finite measurable spaces, morphisms are Markov kernels (stochastic matrices).
- *F* is a contravariant functor from *A* to Kern^f; *F*^a_b : *F*(a) → *F*(b) is denoted element-wise as *F*^a_b(ω_b | ω_a), with ω_b ∈ *F*(b), ω_a ∈ *F*(a).
 - \rightarrow *F* encodes all the ways our data can interact.
 - $\rightarrow \mathscr{A}$ is any poset, not just a collection of subsets.
 - \rightarrow Maps are not just projections.
- $Q = (Q_a \in \mathbb{P}(F(a)), a \in \mathscr{A})$
- *F*_{Bethe}(*Q*) = ∑_{a∈𝔅} *c*(*a*) (𝔼<sub>*Q_a*[*H_a*] − *S*(*Q_a*)); *c*(*a*) = ∑_{b≥a} μ(*b*, *a*) is the generalization of the inclusion-exclusion formula associated with 𝔅.
 </sub>

For a finite poset \mathscr{A} ,

- the 'zeta-operator' of A, denoted ζ, from ⊕_{a∈A} ℝ to ⊕_{a∈A} ℝ is defined as, for any λ ∈ ⊕_{a∈A} ℝ and any a ∈ A, ζ(λ)(a) = ∑_{b≤a} λ_b
- its inverse is denoted as μ ; ($\mu(a, b), b \leq a$) Möbius function of \mathscr{A} .

We want to do Bayesian inference on these diagram.

- Constraint: the Q_a must be compatible under the actions of the F_b^a , i.e. $F_b^a \circ Q_a = Q_b$
- Problem: find an algorithm to 'solve' the optimization problem.

 \rightarrow New message passing algorithm!

F induces several actions: on probabilities, on probabilities seen as vectors, on their dual...

- *˜*F^a_b: ℙ(F(a)) → ℙ(F(b)) is linear map that sends probability distributions p ∈ ℙ(F(a)) to F^a_b ∘ p, we still note *˜*F the linear map from ℝ^{F(a)} to ℝ^{F(b)}.
- *˜F*^{*} is the functor obtained by dualizing the morphisms *˜F*^a_b, i.e.
 ˜F^{*b}_a: *˜F*(*b*)^{*} → *˜F*(*a*)^{*} sends linear maps *I*_b: *˜F*(*b*) → ℝ to
 *I*_b ∘ *˜F*^a_b: *˜F*(*a*) → ℝ.

 μ can be extended to account for \tilde{F} through \tilde{F}^* :

 for a functor *G* from *A* to ℝ-vector spaces, we define μ_G as, for any *a* ∈ *A* and *v* ∈ ⊕_{*a*∈*A*} *G*(*a*), μ_G(*v*)(*a*) = ∑_{*b*≤*a*} μ(*a*, *b*)*G*^{*b*}_{*a*}(*v*_{*b*}).

•
$$\zeta_G$$
 is it's inverse, $\zeta_G(v)(a) = \sum_{b \leq a} G_a^b(v_b)$.

Recall we want to solve $\inf F_{Bethe} = \sum_{a} c(a)F(Q_{a})$ under

- Constraint: the Q_a must be compatible under the actions of the F_b^a , i.e., $F_b^a \circ Q_a = Q_b$
 - i.e., $Q \in \lim \tilde{F}$
 - In fact, no... need to add the condition that the distribution sums to one.
 - But it's okay!

Algorithm 1: Message passage algorithm for presheaves from \mathscr{A} to Kern^f

Data: Initialization: $(m_{a\to b}^0 \in \mathbb{R}^{F(b)}, b, a \in \mathscr{A} \text{ s.t. } b \leq a)$, a poset \mathscr{A} , a presheaf $F : \mathscr{A} \to \operatorname{Kern}^f$; 1 for $t \leq T$ do for $a \in \mathcal{A}, b \in \mathcal{A}$ such that $b \leq a$ do 2 $\forall \omega_a \in F(a), \quad n_{b \to a}(\omega_a) \leftarrow \prod_{\substack{c:b \leq c \\ c \neq a}} \sum_{\omega'_b \in F(b)} m_{c \to b}(\omega'_b) \cdot F^a_b(\omega'_b|\omega_a)$ 3 end 4 for $a \in \mathcal{A}, b \in \mathcal{A}$ such that $b \leq a$ do 5 $b_a = e^{-H_a} \prod_{\substack{b \in \mathscr{A}: \ b \leq a}} n_{b \to a}$ 6 $p_a = rac{b_a}{\sum_{\omega_a} b_a(\omega_a)} p_{a \to b} \leftarrow m_{a \to b} \cdot rac{ ilde{F}_b^a(p_a)}{m_a}$ 7 8 9 end 10 end

• Fix point of this message passing algorithm are critical point of *F*_{Bethe}

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- FE : ∏_{a∈𝒜} P(E_a) → ∏_{a∈𝒜} R defined as FE(Q) = (E_{Qa}[H_a] - S_a(Q_a), a ∈ 𝒜), which sends a collection of probability measures over 𝒜 to their Gibbs free energies.
- $d_Q FE$ denotes the differential of FE at the point Q.

Proposition

Let \mathscr{A} be a finite poset, and let F be a contravariant functor from \mathscr{A} to $Kern^{f}$. Let $H_{a} : F(a) \to \mathbb{R}$ be a collection of (measurable) functions. The critical points of F_{Bethe} are the $Q \in \lim \tilde{F}$ such that:

$$\mu_{\tilde{F}^*} d_Q F E|_{T \lim \tilde{F}} = 0$$

 $T \lim \tilde{F}$ is the underlying vector space of the affine space $\lim \tilde{F}$

Pose $I_a(Q_a) = \mathbb{E}_{Q_a}[H_a] - S(Q_a)$

Theorem (GSP)

F a functor from \mathscr{A}^{op} to vector spaces. An element $u \in \lim \tilde{F}$ is a critical point of the F_{Bethe} if and only if there is $(m_{a \to b} \in \bigoplus_{\substack{a,b:\\b \leq a}} \tilde{F}(b)^*)$ such that for any $a \in \mathscr{A}$,

$$d_{u}l_{a} = \sum_{b \leq a} \tilde{F}_{b}^{a*} \left(\sum_{c \leq b} \tilde{F}_{c}^{b*} m_{b \to c} - \sum_{c \geq b} m_{c \to b} \right)$$
(CP)

- To understand in greater detail these propositions and the previous algorithm, we need to extend the setting of the optimization problem.
- Change the loss:
 - \rightarrow Replace entropy with a "local loss."
 - $\rightarrow S(Q_a) \rightsquigarrow I_a(v_a)$
- Change the functor:
 - \rightarrow Replace *F* with a contravariant functor from a poset \mathscr{A} to Vect.
- Result: we can extend the message passing algorithm to solve:

$$\min_{v}\sum_{a\in\mathscr{A}}c(a)I_a(v_a)$$

with $v := (v_a, a \in \mathscr{A})$ under the constraint $v \in \lim F$.

• This approach is different and on some points more general than decentralized optimization on cellular sheaves [HG19].

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Inference on diagrams

For *F* a functor from \mathscr{A}^{op} to vector spaces, critical points *u* of $\sum_{a \in \mathscr{A}} c(a) l_a(v_a)$ are $u \in \bigoplus_{a \in \mathscr{A}} F(a)$ such that:

$$[\mu_{F^*} d_u I]|_{\lim F} = 0$$

where, $I(v) = (I_a(v_a), a \in \mathscr{A}), d_u I(a) = d_{u_a} I_a$ and,
 $[\mu_{F^*} d_u I](a) = \sum_{b \leq a} \mu(a, b) d_{u_b} I_b \circ F_b^a$

$$0 o \lim F o igoplus_{a \in \mathscr{A}} F(a) \stackrel{\delta_F}{ o} igoplus_{a, b \in \mathscr{A} \atop a \geq b} F(b)$$

where for any $v \in \bigoplus_{\substack{a,b \in \mathscr{A} \\ a \ge b}} F(b)$ and $a, b \in \mathscr{A}$ such that $b \le a$, $\delta_F(v)(a,b) = F_b^a(v_a) - v_b$

This is simply stating that ker $\delta = \lim F$.

Understanding expression of critical points:

$$0 \leftarrow (\lim F)^* \leftarrow \bigoplus_{a \in \mathscr{A}} F(a)^* \xleftarrow{d_F} \bigoplus_{\substack{a,b \in \mathscr{A} \\ a \geq b}} F(b)^*$$

Pose d = δ^* . For any $l_{a \to b} \in \bigoplus_{\substack{a,b \in \mathscr{A} \\ a \geq b}} F(b)^*$ and $a \in \mathscr{A}$,
dm(a) = $\sum F_b^{a^*}(m_{a \to b}) - \sum m_{b \to a}$

a≥b

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b≥a

 $\mu_F^* d_u l \in \operatorname{im} d$

is the same as the fact that there is $(m_{a \rightarrow b} \in F(b)^* | a, b \in \mathscr{A}, b \leq a)$ such that,

 $d_{u}I = \zeta_{F^*} dm$

Assume that the local losses I_a , $a \in \mathscr{A}$ are such that there is a collection of functions g_a , $a \in \mathscr{A}$ that inverses the relation induced by differentiating the local losses, i.e.

$$d_{u_a}l_a = y_a \iff u_a = g_a(y_a)$$

It is the case for the free energy $\mathbb{E}_{Q_a}[H_a] - S(Q_a)$. Messages:

$$m(t) \in \bigoplus_{\substack{a,b:\\b\leq a}} F(b)^*$$
: $m_{a \to b}$ for $b \leq a$

Understanding this choice of message passing algorithm:

g sends Lagrange multipliers *m* to $u \in \bigoplus_{a \in \mathscr{A}} F(a)$. $\delta_F(u) = 0$ defines the constraints on *u*.

 $\delta_F g \zeta_{F^*} d_F$ sends a Lagrange multiplier $m \in \bigoplus_{a,b \in \mathscr{A}} F(b)^*$ to a

constraint $c \in \bigoplus_{\substack{a,b \in \mathscr{A} \\ a \geq b}} F(b)$ defined as, for $a, b \in \mathscr{A}$ such that $b \leq a$,

$$c(a,b) = \delta_F g\zeta_{F^*} \mathsf{d}_F m(a,b) = F_b^a g_a(\zeta_{F^*} \mathsf{d}_F m(a)) - g_b(\zeta_{F^*} \mathsf{d}_F m(b)))$$
(0.2)

We are interested in c = 0, i.e.

$$\delta_F g \zeta_{F^*} \mathsf{d}_F m = 0$$

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Understanding this choice of message passing algorithm:

Choice of algorithm on the Lagrange multipliers so that $\delta_F g \zeta_{F^*} d_F m = 0$,

$$m(t+1) - m(t) = \delta_F g \zeta_{F^*} d_F m(t)$$

Any other choice would also be a good candidate!

The message passing algorithm is defined as:

$$\delta m := \delta_{\tilde{F}} g \zeta_{\tilde{F}^*} d_{\tilde{F}} m$$

Define $BP_{F,H} := \delta_{\tilde{F}} g \zeta_{\tilde{F}^*} d_{\tilde{F}}$. When differentiating the free energy:

$$y_a = H_a + \ln q_a + 1$$

Therefore, $g_a(y_a) = e^{y_a - H_a - 1}$.

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- Joint work with Toby St Clere Smithe, in progress.
- Consider a natural transformation φ : F → F₁ where φ_a is a deterministic map, not a Markov kernel.
- The map ϕ extends into maps between $\tilde{F} \to \tilde{F}_1$ and $\tilde{F}^* \to \tilde{F}_1^*$.
- ϕ induces maps between $\bigoplus_b \tilde{F}(b) \to \bigoplus_b \tilde{F}_1(b)$ and $\bigoplus_{b \leq a} \tilde{F}(b) \to \bigoplus_{b \leq a} \tilde{F}_1(b)$. It also induces a map $\phi^* : \bigoplus_b \tilde{F}_1^*(b) \to \bigoplus_b \tilde{F}^*(b)$ and $\phi^* : \bigoplus_{b \leq a} \tilde{F}_1^*(b) \to \bigoplus_{b \leq a} \tilde{F}^*(b)$.

Pose:

$$ilde{\mathcal{H}}_{a} = \ln \sum_{\omega': \phi_{a}(\omega') = \omega} e^{-\mathcal{H}_{a}(\omega')}$$

Then we showed that:

$$BP_{F_1,\tilde{H}} = \phi \circ BP_{F,H} \circ \phi^*$$

- $\rightarrow\,$ Few results on characterizing critical points of the Bethe free energy.
- $\rightarrow\,$ Use transformations on the underlying functor to reduce to simpler cases (Hamiltonians, posets).
- What about base change? $\phi : \mathscr{A} \to \mathscr{A}_1$
 - $\rightarrow\,$ When a right adjoint to the pullback exists, results on natural transformations can be reused.
 - $\rightarrow\,$ When \mathscr{A} is isomorphic to a full subposet of $\mathscr{A}_1,$ similar result holds.

Thank you for your attention!

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Grégoire Sergeant-Perthuis,

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